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Bad(s, t) IS HYPERPLANE ABSOLUTE WINNING

EREZ NESHARIM AND DAVID SIMMONS

ABSTRACT. J. An [A] proved that for any $s, t \geq 0$ such that $s + t = 1$, **Bad**(s, t) is $(34\sqrt{2})^{-1}$ -winning for Schmidt's game. We show that using the main lemma from [A] one can derive a stronger result, namely that **Bad**(s, t) is hyperplane absolute winning in the sense of [BFKRW]. As a consequence one can deduce the full dimension of **Bad**(s, t) intersected with certain fractals.

1. STATEMENT OF RESULTS

Throughout this paper, fix $s, t \geq 0$ with $s + t = 1$. Let **Bad**(s, t) denote the set

$$\mathbf{Bad}(s, t) = \{(x, y) \in \mathbb{R}^2 : \inf_{q \in \mathbb{N}} \max(q^s \|qx\|, q^t \|qy\|) > 0\},$$

where $\|\cdot\|$ is the distance to the nearest integer. Schmidt's conjecture, proven in [BPV], states that $\mathbf{Bad}(\frac{1}{3}, \frac{2}{3}) \cap \mathbf{Bad}(\frac{2}{3}, \frac{1}{3})$ is nonempty. A stronger result was proven by J. An [A, Theorem 1.1], who showed that **Bad**(s, t) is $(34\sqrt{2})^{-1}$ -winning for Schmidt's game. In particular this implies (cf. [S, Theorem 2] and [S, Corollary 2 of Theorem 6]) that for any countable collection of pairs $(s_n, t_n)_{n=1}^\infty$, the intersection $\bigcap_{n \in \mathbb{N}} \mathbf{Bad}(s_n, t_n)$ is nonempty and in fact has full Hausdorff dimension in \mathbb{R}^2 .

The object of this note is to give a proof of the following strengthening of An's theorem:

Theorem 1.1. *The set **Bad**(s, t) is hyperplane absolute winning in the sense of [BFKRW].*

Theorem 1.1 is a generalization of An's theorem since every hyperplane absolute winning set is α -winning for Schmidt's game for every $0 < \alpha < 1/2$ [BFKRW, Proposition 2.3(a)]. Moreover, the intersection of a hyperplane absolute winning set with a hyperplane diffuse set which is the support of an Ahlfors regular measure (see [BFKRW] for the definitions) has full dimension with respect to that set [BFKRW, Theorems 4.7 and 5.3]. In particular hyperplane absolute winning sets have full dimension intersection with many well-known fractals such as the Sierpinski triangle and the von Koch snowflake curve. Finally, the class of hyperplane absolute winning sets is closed under countable intersections [BFKRW, Proposition 2.3(b)], and invariant under \mathcal{C}^1 diffeomorphisms [BFKRW, Proposition 2.3(c)]. As a result we have the following:

Corollary 1.2. *For any hyperplane diffuse set $K \subseteq \mathbb{R}^2$ which is the support of an Ahlfors regular measure, for any countable collection of pairs $(s_n, t_n)_{n=1}^\infty$, and for any countable collection of \mathcal{C}^1 diffeomorphisms $(f_n)_{n=1}^\infty$ from \mathbb{R}^2 to itself, the intersection*

$$K \cap \bigcap_{n \in \mathbb{N}} f_n(\mathbf{Bad}(s_n, t_n))$$

has full dimension in K .

We remark also that while the strategy for Schmidt's game given in An's paper is not explicit, depending on König's lemma (cf. [A, Proposition 2.2]), the strategy which we give in the proof of Theorem 1.1 can be described more or less explicitly; see Remark 2.5.

Remark 1.3. The cases $s = 0$ and $t = 0$ of Theorem 1.1 are trivial consequences of the fact that the set of badly approximable numbers is absolute winning (see [M, Theorem 1.3] or [BFKRW, Theorem 2.5]) and will be omitted. Throughout the proof we assume that $s, t > 0$.

Remark 1.4. Although the higher-dimensional analogue of Schmidt's conjecture has been established by V. V. Beresnevich [B], it is still not known, for example, whether **Bad**(s, t, u) is winning for all $s, t, u \geq 0$ with $s + t + u = 1$, where **Bad**(s, t, u) is defined appropriately.

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2. PRELIMINARIES

The proof of Theorem 1.1 will consist of combining the main idea of [A] with the main idea of [FSU, Appendix C]. We therefore begin by recalling these ideas.

2.1. The main lemma of [A]. For each $P = \left(\frac{p}{q}, \frac{r}{q}\right) \in \mathbb{Q}^2$ and $\varepsilon > 0$, following [A] we let¹

$$\Delta_\varepsilon(P) = \left\{ (x, y) \in \mathbb{R}^2 : \left| x - \frac{p}{q} \right| \leq \frac{\varepsilon}{q^{1+s}} \text{ and } \left| y - \frac{r}{q} \right| \leq \frac{\varepsilon}{q^{1+t}} \right\},$$

so that

$$(2.1) \quad \mathbf{Bad}(s, t) = \bigcup_{\varepsilon > 0} \left(\mathbb{R}^2 \setminus \bigcup_{P \in \mathbb{Q}^2} \Delta_\varepsilon(P) \right).$$

Let \mathcal{L} denote the collection of lines (hyperplanes) in \mathbb{R}^2 . If $L \in \mathcal{L}$ and $\gamma > 0$, we let $L^{(\gamma)}$ denote the γ -thickening of L , i.e. the set $\{(x, y) \in \mathbb{R}^2 : d((x, y), L) \leq \gamma\}$.

Lemma 2.1 ([A, Lemma 4.2]). *Fix $R > 1$ and $\ell > 0$. There exists $\varepsilon > 0$ and a partition*

$$(2.2) \quad \mathbb{Q}^2 = \bigcup_{\delta=1}^2 \bigcup_{n=1}^{\infty} \bigcup_{k=1}^n \mathcal{P}_{n,k}^{(\delta)}$$

such that the following holds: For each $m \geq 0$, let

$$\mathcal{P}_m = \bigcup_{\delta=1}^2 \bigcup_{k=1}^m \mathcal{P}_{m,k}^{(\delta)},$$

and let \mathcal{B}_m denote the collection of balls $B \subseteq \mathbb{R}^2$ of radius $R^{-m}\ell/2$ satisfying

$$(2.3) \quad \forall m' \leq m \quad \forall P \in \mathcal{P}_{m'}, \quad \Delta_\varepsilon(P) \cap B = \emptyset.$$

Then for all $n \geq k \geq 1$, for all $\delta \in \{1, 2\}$, and for all $B \in \mathcal{B}_{n-k}$, there is a line $L = L_{n,k,\delta}(B) \in \mathcal{L}$ such that

$$(2.4) \quad \Delta_\varepsilon(P) \subseteq L^{(\frac{1}{3}\ell R^{-n})} \quad \forall P \in \mathcal{P}_{n,k}^{(\delta)} \text{ such that } \Delta_\varepsilon(P) \cap B \neq \emptyset.$$

Remark 2.2. The relation between Lemma 2.1 and [A] requires some explanation. First of all, given $R > 1$ and $\ell > 0$, one can let $\varepsilon > 0$ be defined by the equation [A, (3.2)]. Next, one can define the partition (2.2) as in [A, p. 5-6]. At this point [A, Lemma 4.2] can almost be read as stated, except that An has fixed $\tau \in \mathcal{S}_{n-k}$ instead of $B \in \mathcal{B}_{n-k}$, and has considered the set $\mathcal{P}_{n,k}^{(\delta)}(\tau) = \{P \in \mathcal{P}_{n,k}^{(\delta)} : \Phi(\tau) \cap \Delta_\varepsilon(P) \neq \emptyset\}$ in place of the set $\{P \in \mathcal{P}_{n,k}^{(\delta)} : \Delta_\varepsilon(P) \cap B \neq \emptyset\}$. But we observe that for $\tau \in \mathcal{T}_{n-k}$, we have $\tau \in \mathcal{S}_{n-k}$ if and only if $\Phi(\tau) \in \mathcal{B}_{n-k}$.² Moreover, the proof of [A, Lemma 4.2] works equally well if $\Phi(\tau)$ is replaced by an arbitrary element in \mathcal{B}_{n-k} . Thus the lemma holds just as well if $\Phi(\tau)$ denotes an arbitrary element of \mathcal{B}_{n-k} rather than an arbitrary element of $\Phi(\mathcal{S}_{n-k})$.

¹We remark that the c in [A] corresponds to our ε ; our c corresponds to the c in [FSU, Appendix C].

²Here we ignore the distinction between balls and squares. The difference is important only in calculating diameter; the diameter of a square with respect to the max norm is equal to its side length, while the diameter of a ball is equal to twice its radius. This is why we require balls in \mathcal{B}_m to have radius $R^{-m}\ell/2$, while if $\tau \in \mathcal{S}_m$, the side length of $\Phi(\tau)$ is $R^{-m}\ell$.

2.2. Two variants of Schmidt's game. We proceed to describe two variants of Schmidt's game, one introduced in [BFKRW] and the other introduced in [FSU, Appendix C]. In this paper we will not deal directly with the first game, but we will prove that **Bad**(s, t) is winning with respect to the second game. Since the two games are equivalent (Lemma 2.4 below), this proves that **Bad**(s, t) is also winning with respect to the first game, and therefore has the large dimension properties described in the introduction.

Definition 2.3. Given $0 < \beta < 1/3$, Alice and Bob play the β -hyperplane absolute game as follows:

1. Bob begins by choosing a ball $B_0 = B(z_0, r_0) \subseteq \mathbb{R}^2$.
2. On Alice's n th turn, she chooses a set of the form $L_n^{(\tilde{r}_n)}$ with $L_n \in \mathcal{L}$, $0 < \tilde{r}_n \leq \beta r_n$, where r_n is the radius of Bob's n th move $B_n = B(z_n, r_n)$. We say that Alice *deletes* her choice $L_n^{(\tilde{r}_n)}$.
3. On Bob's $(n + 1)$ st turn, he chooses a ball $B_{n+1} = B(z_{n+1}, r_{n+1})$ satisfying

$$(2.5) \quad r_{n+1} \geq \beta r_n \text{ and } B_{n+1} \subseteq B_n \setminus L_n^{(\tilde{r}_n)},$$

where $B_n = B(z_n, r_n)$ was his n th move, and $L_n^{(\tilde{r}_n)}$ was Alice's n th move.

4. If $r_n \not\rightarrow 0$, then Alice wins by default. Otherwise, the balls $(B_n)_1^\infty$ intersect at a unique point which we call the *outcome* of the game.

If Alice has a strategy guaranteeing that the outcome lies in a set S (or that she wins by default), then the set S is called β -hyperplane absolute winning. If a set S is β -hyperplane absolute winning for all $0 < \beta < 1/3$ then it is called *hyperplane absolute winning*.

By contrast, given $\beta, c > 0$, Alice and Bob play the (β, c) -hyperplane potential game as follows:

1. Bob begins by choosing a ball $B(\mathbf{x}_0, r_0) \subseteq \mathbb{R}^2$.
2. For each n , after Bob makes his n th move $B_n = B(\mathbf{x}_n, r_n)$, Alice will make her n th move. She does this by choosing a countable collection of sets of the form $L_{i,n}^{(r_{i,n})}$, with $L_{i,n} \in \mathcal{L}$ and $r_{i,n} > 0$, satisfying

$$(2.6) \quad \sum_i r_{i,n}^c \leq (\beta r_n)^c.$$

3. After Alice makes her n th move, Bob will make his $(n + 1)$ st move by choosing a ball $B_{n+1} = B(\mathbf{x}_{n+1}, r_{n+1})$ satisfying

$$(2.7) \quad r_{n+1} \geq \beta r_n \text{ and } B_{n+1} \subseteq B_n,$$

where $B_n = B(\mathbf{x}_n, r_n)$ was his n th move.

4. If $r_n \not\rightarrow 0$, then Alice wins by default. Otherwise, the balls $(B_n)_1^\infty$ intersect at a unique point which we call the *outcome* of the game. If the outcome is an element of any of the sets $L_{i,n}^{(r_{i,n})}$ which Alice chose during the course of the game, she wins by default.

If Alice has a strategy guaranteeing that the outcome lies in a set S (or that she wins by default), then the set S is called (β, c) -hyperplane potential winning. If a set is (β, c) -hyperplane potential winning for all $\beta, c > 0$, then it is *hyperplane potential winning*.

The following lemma is a special case of the main result of [FSU, Appendix C]:

Lemma 2.4 ([FSU, Theorem C.8]). *A set is hyperplane potential winning if and only if it is hyperplane absolute winning.*

Sketch of the proof. We sketch only the forward direction, as it is the one which we use. Suppose that $S \subseteq \mathbb{R}^d$ is hyperplane potential winning. Let $\beta > 0$. Fix $\tilde{\beta}, c > 0$ small to be determined, and consider a strategy of Alice which is winning for the $(\tilde{\beta}, c)$ -hyperplane potential game. Each time Bob makes a move $B_n = B(z_n, r_n)$, Alice chooses a collection of sets $\{L_{i,n}^{(r_{i,n})}\}_{i=1}^{N_n}$ (with $N_n \in \mathbb{N} \cup \{\infty\}$) satisfying (2.6). Alice's corresponding strategy in the β -hyperplane absolute game will be to choose her set $L_n^{(\beta r_n)} \subseteq \mathbb{R}^2$ so as to

maximize

$$(2.8) \quad \phi(B_n; L_n^{(\beta r_n)}) := \sum_{m=0}^n \sum_{\substack{i \\ L_{i,m}^{(r_{i,m})} \subseteq L_n^{(\beta r_n)} \\ L_{i,m}^{(r_{i,m})} \cap B_n \neq \emptyset}} r_{i,m}^c.$$

Suppose that Alice plays according to this strategy, and let $(B_n)_1^\infty$ be the sequence of Bob's moves. For each $n \in \mathbb{N}$ let

$$\phi(B_n) = \sum_{m=0}^n \sum_{L_{i,m}^{(r_{i,m})} \cap B_n \neq \emptyset} r_{i,m}^c.$$

One demonstrates by induction on n (see [FSU, Appendix C] for details) that if $\tilde{\beta}$ and c are chosen sufficiently small, then there exists $\varepsilon > 0$ such that for all $n \in \mathbb{N}$,

$$(2.9) \quad \phi(B_n) \leq (\varepsilon r_n)^c.$$

Intuitively, the reason for this is that Alice is “deleting the regions with high ϕ -value”, and is therefore minimizing the ϕ -value of Bob's balls. Thus she is forcing $\phi(B_n)$ to be as small as possible.

Now suppose that $r_n \rightarrow 0$; otherwise Alice wins the β -hyperplane absolute game by default. Then (2.9) implies that $\phi(B_n) \rightarrow 0$. In particular, for each (i, m) , $r_{i,m}^c > \phi(B_n)$ for all n sufficiently large which implies $L_{i,m}^{(r_{i,m})} \cap B_n = \emptyset$. Thus the outcome of the game does not lie in $L_{i,m}^{(r_{i,m})}$ for any (i, m) , so Alice does not win the $(\tilde{\beta}, c)$ -hyperplane potential game by default. Thus if she wins, then she must win by having the outcome lie in S . Since the outcome is the same for the $(\tilde{\beta}, c)$ -hyperplane potential game and the β -hyperplane absolute game, this implies that she also wins the β -hyperplane absolute game. \square

Remark 2.5. The proof of Lemma 2.4 shows that if Alice has an explicit strategy to win the hyperplane potential game, then she also has an explicit strategy to win the hyperplane absolute game. Here by “explicit” we mean roughly that Alice can calculate what her next move should be in a finite amount of time, based on the input of Bob's previous move. To see that Alice's strategy in the above proof can be made explicit, note that Alice does not actually have to maximize (2.8) precisely; for example, if she instead makes a choice for which the value of (2.8) is at least $1/2$ of the maximum value, then the proof will still go through with minor modifications. This margin of error allows her to restrict her attention to a finite set of potential moves; moreover, for each potential move she can truncate the series (2.8) at an appropriate point to make it a finite series. Thus she can compute an “approximate maximum” explicitly.

3. PROOF OF THEOREM 1.1

In this section we prove Theorem 1.1, which states that $\mathbf{Bad}(s, t)$ is hyperplane absolute winning.

Proof of Theorem 1.1. By Lemma 2.4, it suffices to show that $\mathbf{Bad}(s, t)$ is hyperplane potential winning. Fix $\beta, c > 0$, and we will show that it is (β, c) -hyperplane potential winning. Let $B(\mathbf{x}_0, r_0) \subseteq \mathbb{R}^2$ be Bob's first move. Fix $R \geq \beta^{-1}$ to be determined (depending on β and c), and let $\ell = 2r_0$. Let $\varepsilon > 0$ and the partition (2.2) be as in Lemma 2.1. Alice's strategy will be defined by infinitely many “triggers” as follows: For each $m \geq 0$, Alice will wait until Bob chooses a ball $B_j = B(\mathbf{x}_j, r_j)$ that satisfies $r_j \leq R^{-m}r_0/2$. The first j for which this inequality holds will be denoted j_m , with $j_m = \infty$ if it never holds. We observe that

- (i) $j_m \geq 1$ for all $m \geq 0$, since $r_0 > R^{-m}r_0/2$, and
- (ii) $j_{m+1} \geq j_m + 1$, since $r_{j_m} \geq \beta r_{j_m-1} > \beta R^{-m}r_0/2 \geq R^{-(m+1)}r_0/2$ (using the fact that $R \geq \beta^{-1}$).

Fix $m \geq 0$, and let $j = j_m$. Let $\tilde{B}_j = B(\mathbf{x}_j, R^{-m}r_0)$. On Alice's j th turn, her strategy will be as follows:

- (1) If $\tilde{B}_j \notin \mathcal{B}_m$, then she will do nothing.
- (2) If $\tilde{B}_j \in \mathcal{B}_m$, then for each $k \geq 1$ and $\delta \in \{1, 2\}$ she will apply Lemma 2.1 to get a line $L_{k,\delta} = L_{m+k,k,\delta}(\tilde{B}_j) \in \mathcal{L}$, and she will delete the hyperplane-neighborhood $L_{k,\delta}^{(3R^{-(m+k)}r_0)}$.

The legality of this action is guaranteed by (ii), which shows that Alice is not deleting multiple collections on the same turn, together with the inequality

$$\begin{aligned} \sum_{\delta=1}^2 \sum_{k=1}^{\infty} \left(3R^{-(m+k)}r_0\right)^c &= (3R^{-m}r_0)^c 2 \sum_{k=1}^{\infty} R^{-ck} \\ &\leq (\beta^2 R^{-m}r_0/2)^c && \text{(for } R \text{ chosen large enough)} \\ &< (\beta r_{j_m})^c. && \text{(since } r_{j_m} \geq \beta r_{j_m-1} > \beta R^{-m}r_0/2) \end{aligned}$$

To complete the proof, we must show that the strategy described above guarantees a win for Alice. By contradiction, suppose that Bob can play in a way so that Alice loses. By definition, this means that the radii of Bob's balls tend to zero, and that their intersection point $\mathbf{x} \in \mathbb{R}^2$ is not in **Bad**(s, t) nor in any of the hyperplane-neighborhoods which Alice deleted in the course of the game. In particular, the radii tending to zero means that each of the triggers happens eventually, i.e. $j_m < \infty$ for all $m \geq 0$.

Claim 3.1. For all $m \geq 0$, $\tilde{B}_{j_m} := B(\mathbf{x}_{j_m}, R^{-m}r_0) \in \mathcal{B}_m$.

Proof. We proceed by strong induction and contradiction. Suppose the claim holds for all $0 \leq m < M$, but does not hold for M . Then there exist $M' \leq M$ and $P \in \mathcal{P}_{M'}$ such that $\Delta_\varepsilon(P) \cap \tilde{B}_J \neq \emptyset$, where $J = j_{M'}$. Write $P \in \mathcal{P}_{M',k}^{(\delta)}$ for some $1 \leq k \leq M'$ and $\delta \in \{1, 2\}$. Let $m = M' - k < M$, and let $j = j_m$. We apply the induction hypothesis to see that $\tilde{B}_j \in \mathcal{B}_m$. Thus on Alice's j th turn, she must have deleted the hyperplane-neighborhood $A := L_{k,\delta}^{(3R^{-(m+k)}r_0)}$, where $L_{k,\delta} = L_{m+k,k,\delta}(\tilde{B}_j)$ is as in Lemma 2.1.

On the other hand, since $J \geq j$, we have $B_J \subseteq B_j$; thus $d(\mathbf{x}_j, \mathbf{x}_J) \leq r_j \leq R^{-m}r_0/2$ and so $\tilde{B}_J \subseteq \tilde{B}_j$. Combining this with the contradiction hypothesis gives $\Delta_\varepsilon(P) \cap \tilde{B}_j \neq \emptyset$. So by the definition of $L_{k,\delta} = L_{m+k,k,\delta}(\tilde{B}_j)$ (cf. (2.4)), we have

$$\Delta_\varepsilon(P) \subseteq L_{k,\delta}^{\left(\frac{2}{3}R^{-(m+k)}r_0\right)}.$$

Since $\frac{2}{3}R^{-(m+k)}r_0 + 2R^{-M}r_0 \leq 3R^{-(m+k)}r_0$, this implies $(\Delta_\varepsilon(P))^{(2R^{-M}r_0)} \subseteq A$. In particular, since $\Delta_\varepsilon(P) \cap \tilde{B}_J \neq \emptyset$ we have

$$\mathbf{x} \in B_J \subseteq \tilde{B}_J \subseteq (\Delta_\varepsilon(P))^{(2R^{-M}r_0)} \subseteq A.$$

This demonstrates that Alice won by default, contradicting our hypothesis. ◁

Now for all $P \in \mathbb{Q}^2$, we have $P \in \mathcal{P}_m$ for some $m \geq 1$. Let $j = j_m$. Applying Claim 3.1, we see that $\Delta_\varepsilon(P) \cap \tilde{B}_j = \emptyset$. But $\mathbf{x} \in B_j \subseteq \tilde{B}_j$, so $\mathbf{x} \notin \Delta_\varepsilon(P)$. By (2.1), this means $\mathbf{x} \in \mathbf{Bad}(s, t)$. So Alice won, contradicting our hypothesis. This completes the proof. ◻

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