

Cutoff for a Random Walk on the Integers mod n

Michael E. Bate and Stephen B. Connor

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Abstract

We analyse a random walk on the ring of integers mod n , which at each time point can make an additive ‘step’ or a multiplicative ‘jump’. When the probability of making a jump tends to zero as an appropriate power of n we prove the existence of a total variation cutoff for this process, with cutoff time dependent on whether the step distribution has zero mean.

Keywords: random walk; mixing time; cutoff phenomenon; group representation theory; integers mod n ; random number generation

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1 Introduction

In this note we consider a random walk X on $\mathbb{Z}_n = \mathbb{Z}/n\mathbb{Z}$ (where n is *odd*) defined as follows:

$$X_t = \begin{cases} X_{t-1} + \xi'_t \pmod n & \text{with probability } 1 - p_n \\ 2X_{t-1} \pmod n & \text{with probability } p_n, \end{cases} \quad (1)$$

where $\{\xi'_t\}$ are a set of i.i.d. random variables with finite support $B \subset \mathbb{Z}$, whose distribution does not vary with n . We denote the mean and variance of ξ' by μ and σ^2 respectively. We will refer to an ‘addition’ move as a ‘step’, and to a ‘multiplication’ move as a ‘jump’. To ensure that X is irreducible we assume that the group $\langle B_n, + \rangle$ is not a proper subgroup of \mathbb{Z}_n for any odd n , where $B_n = \{z \pmod n : z \in B\}$. Furthermore, since n is odd, multiplication by 2 is an invertible operation, and thus X is ergodic with uniform equilibrium distribution π_n on \mathbb{Z}_n .

Define the total variation distance from π_n of a probability distribution P on \mathbb{Z}_n by

$$\|P - \pi_n\|_{\text{TV}} = \max_{A \subset \mathbb{Z}_n} |P(A) - \pi_n(A)| = \frac{1}{2} \sum_{s \in \mathbb{Z}_n} |P(s) - 1/n|.$$

A number of authors have previously considered random processes of the form

$$X_t = a_t X_{t-1} + b_t \pmod n;$$

these processes are similar to schemes used for random number generation, a link which has naturally motivated interest in bounding the time taken for the total variation distance from uniform to become suitably small (the so-called ‘mixing time’, which is typically taken to

be the first time at which the total variation distance drops below $1/4$). A nice introduction to the area can be found in [Terras \(1999, Chapter 6\)](#). The earliest such work appears to be that of [Chung, Diaconis, and Graham \(1987\)](#), in which $a_t = a = 2$ and b_t is chosen uniformly from $\{-1, 0, 1\}$: they show that $O(\log n \log \log n)$ steps suffice for this walk to mix, and that $O(\log n \log \log n)$ steps are also necessary for n of the form $2^m - 1$; on the other hand, for almost all odd n , $1.02 \log_2 n$ steps suffice. This (deterministic) act of doubling each time causes the process to mix significantly faster than when $a_t = 1$ for all t where, if b_t is uniform on a finite set (and assuming that the resulting process is irreducible), the mixing time is of order n^2 ([Diaconis, 1988](#); [Saloff-Coste, 2004](#)).

Rather more general results have been established in a series of works by Hildebrand. It is shown in his thesis ([Hildebrand, 1992, Chapter 3](#)) that if multiplication is deterministic ($a_t = a$ for all t) and for fairly general choices of b_t (which don't depend on n), $O(\log n \log \log n)$ steps suffice, and in fact for almost all n , $O(\log n)$ steps suffice; the method closely follows that of [Chung et al. \(1987\)](#). When a_t is allowed to vary with t , a general upper bound for the mixing time is proved in ([Hildebrand, 1993](#)): using a recursive relation involving discrete Fourier transforms (of which more below), he shows that (unless $a_t = 1$ always, $b_t = 0$ always, or a_t and b_t can each take on only one value) $O((\log n)^2)$ time steps are always sufficient. Other related results can be found in [Hildebrand \(1994a,b\)](#).

A particularly interesting feature of these processes is the quantitatively different behaviour that can be obtained by making small changes to the distribution of a_t and b_t . For example, [Chung et al. \(1987\)](#) remark upon the following curiosity to be found when $a_t = 2$ and b_t is supported on $\{-1, 0, 1\}$ with $\mathbb{P}(b_t = 1) = \mathbb{P}(b_t = -1) = q$. If $q = 1/4$ or $q = 1/2$ then $O(\log n)$ steps suffice to make the total variation distance small; however, if $q = 1/3$ then $O(\log n \log \log n)$ steps may be required. Similarly, [Hildebrand \(1992, Chapter 5\)](#) considers the situation where b_t is uniform on ± 1 and a_t is supported on $\{2, (n+1)/2\}$, with $\mathbb{P}(a_t = 2) = p \in (0, 1)$: the mixing time is shown to be at most $O((\log n)^m)$, where m is 2 if $p = 1/2$, and 1 otherwise. If the distribution of b_t is altered to uniform on $\{-1, 0, 1\}$ then $O((\log n \log \log n)^m)$ steps suffice.

The principal difference between these earlier works and the process defined in (1) is that we allow the probability of a ‘jump’, p_n , to depend on n . In particular, we are able to show that if p_n tends to zero as a power of n , then our process exhibits a total variation cutoff.

Definition 1. ([Levin, Peres, and Wilmer, 2009](#)) A sequence of Markov chains $\{X^{(n)}\}_{n \in \mathbb{N}}$ is said to exhibit a *total variation cutoff* at time τ_n with *window size* w_n if $w_n = o(\tau_n)$ and

$$\begin{aligned} \lim_{c \rightarrow \infty} \liminf_{n \rightarrow \infty} \|\mathbb{P}(X_{\tau_n - cw_n}^{(n)} \in (\cdot)) - \pi_n(\cdot)\|_{\text{TV}} &= 1 \\ \lim_{c \rightarrow \infty} \limsup_{n \rightarrow \infty} \|\mathbb{P}(X_{\tau_n + cw_n}^{(n)} \in (\cdot)) - \pi_n(\cdot)\|_{\text{TV}} &= 0. \end{aligned}$$

Intuitively this says that as n gets large the convergence to equilibrium, measured using total variation distance, happens in a negligible window of order w_n around the cutoff time τ_n . We remark that it is possible for the ‘right’ and ‘left’ window sizes in the above definition to be of different orders – see [Connor \(2010\)](#) for an example. There has been much interest in studying the mixing times of Markov chains and proving the existence of cutoff phenomena: see [Levin et al. \(2009\)](#) and [Diaconis \(2011\)](#) for recent introductions to the area, or [Saloff-Coste \(2004\)](#) for a more analytical overview.

Our main result is the following.

Theorem 2. *Suppose that $p_n = 1/(2n^\alpha)$ for some $\alpha > 0$ such that $n/\sigma_{S'} \rightarrow \infty$, where*

$$\sigma_{S'}^2 := \frac{(1 - p_n)(\mu^2 + p_n\sigma^2)}{p_n^2}.$$

Then X exhibits a total variation cutoff at time $T_n^X = 2n^\alpha \log_2(n/\sigma_{S'})$, with window size $n^{\alpha/2}\sqrt{T_n^X}$.

This paper's contribution contrasts with the existing results mentioned above for processes of the type $X_t = a_t X_{t-1} + b_t$ (where the distribution of a_t is independent of n), for which (to the best of our knowledge) no cutoff results have been established. In the present setting the mixing time of our process X is also relatively insensitive to the distribution of the step lengths ξ'_t . Theorem 2 shows that the mixing time of X essentially depends on ξ' only through its mean, μ : in the case of zero drift the mixing time is $2(1 - \alpha/2)n^\alpha \log_2 n$ (for $0 < \alpha < 2$), while if $\mu \neq 0$ the chain mixes slightly faster, with cutoff at $2(1 - \alpha)n^\alpha \log_2 n$ (for $0 < \alpha < 1$).

2 Working with a subsampled chain

The main obstruction to analysing our process X using standard techniques for random walks on groups is that the distribution of X_k is not given by convolution of k independent increment distributions. This problem can be overcome by (initially) restricting attention to the process Y which is produced by subsampling X at jump times. Denote the jump times of X by τ_1, τ_2, \dots , and let $\tau_0 = 0$; then $Y_k := X_{\tau_k}$. This process clearly satisfies $Y_k = Y'_k \pmod n$, where

$$Y'_k = 2^k Y_0 + \sum_{i=1}^k 2^{k+1-i} S'_i, \quad (2)$$

and where

$$S'_i = \sum_{t=\tau_{i-1}+1}^{\tau_i-1} \xi'_t. \quad (3)$$

Here (and throughout) we use the convention that random variables with a prime take values in \mathbb{Z} , while those without take values in \mathbb{Z}_n . Thus $S_i = S'_i \pmod n$ is the change in X due to steps taken between jump times τ_{i-1} and τ_i . Like X , Y is ergodic with uniform equilibrium distribution. From (2) it is clear that the mixing time of Y is independent of its starting state, and so for ease of exposition we shall set $X_0 = Y_0 = 0$. It is also clear from (2) that the distribution of Y_k is given by convolution of the distributions corresponding to the independent increments $\{2^{k+1-i} S'_i\}$, and this will prove essential to our method for establishing an upper bound on the mixing time in Section 3.

The length of time between jumps of X clearly has a Geometric(p_n) distribution:

$$\mathbb{P}(\tau_1 = j) = p_n(1 - p_n)^{j-1}, \quad j = 1, 2, \dots,$$

and a straightforward application of the conditional variance formula shows that

$$\sigma_{S'}^2 := \text{Var}[S'_i] = \frac{(1 - p_n)(\mu^2 + p_n\sigma^2)}{p_n^2}. \quad (4)$$

Note in particular that when $p_n \rightarrow 0$,

$$\sigma_{S'}^2 \sim \begin{cases} \sigma^2/p_n & \text{if } \mu = 0 \\ \mu^2/p_n^2 & \text{otherwise.} \end{cases} \quad (5)$$

Theorem 3. *Suppose that $p_n = 1/(2n^\alpha)$ for some $\alpha > 0$ such that $n/\sigma_{S'} \rightarrow \infty$ as $n \rightarrow \infty$. Then Y exhibits a cutoff (in total variation distance) at time $T_n = \log_2(n/\sigma_{S'})$, with cutoff window of size $O(1)$. Indeed,*

$$\|\mathbb{P}(Y_{T_n+c} \in \cdot) - \pi_n(\cdot)\|_{\text{TV}} \begin{cases} \geq 1 - 4^{1+c/3} & c < 0 \\ \leq O(4^{-c}) & \text{as } c \rightarrow \infty. \end{cases}$$

The fact that Y exhibits such a tight cutoff makes it simple to demonstrate a cutoff for X , as claimed in Theorem 2.

Corollary 4. *In the setting of Theorem 3, X exhibits a cutoff at time $T_n^X := T_n/p_n$, with window size $\sqrt{T_n}/p_n$.*

Proof. Let $w_n = \sqrt{T_n}/p_n$, and for $c \in \mathbb{R}$ let $J_n(c)$ denote the number of jumps in X before time $T_n^X + cw_n$. Then $J_n(c) \sim \text{Poisson}(m_n(c))$, where $m_n(c) = p_n(T_n^X + cw_n) = T_n + c\sqrt{T_n}$. The proof essentially now follows from the observation that $J_n(c)$ concentrates in an interval of order $\sqrt{T_n}$ around $m_n(c)$ as $n \rightarrow \infty$. Indeed,

$$\limsup_{n \rightarrow \infty} \mathbb{P}\left(|J_n(c) - m_n(c)| > |c| \sqrt{T_n}\right) \leq \limsup_{n \rightarrow \infty} \frac{\text{Var}[J_n(c)]}{c^2 T_n} = \frac{1}{c^2}.$$

To show that X has not mixed before time T_n^X we simply note that, since Y exhibits a cutoff at T_n with window size $O(1)$, for $c \geq 0$

$$\begin{aligned} \liminf_{n \rightarrow \infty} \|\mathbb{P}(X_{T_n^X - cw_n} \in \cdot) - \pi_n(\cdot)\|_{\text{TV}} &\geq 1 - \limsup_{n \rightarrow \infty} \mathbb{P}(J_n(-c) \geq T_n - O(1)) \\ &\geq 1 - \limsup_{n \rightarrow \infty} \mathbb{P}\left(|J_n(-c) - m_n(-c)| > |c| \sqrt{T_n}\right) \geq 1 - \frac{1}{c^2}. \end{aligned}$$

Similarly,

$$\begin{aligned} \limsup_{n \rightarrow \infty} \|\mathbb{P}(X_{T_n^X + cw_n} \in \cdot) - \pi_n(\cdot)\|_{\text{TV}} &\leq \limsup_{n \rightarrow \infty} \|\mathbb{P}(Y_{J_n(c)} \in \cdot) - \pi_n(\cdot)\|_{\text{TV}} \\ &\leq \limsup_{n \rightarrow \infty} \mathbb{P}(J_n(c) \leq T_n + O(1)) \leq \frac{1}{c^2}. \end{aligned}$$

□

It therefore remains to prove Theorem 3. The left hand cutoff window follows relatively simply from an application of Chebychev's inequality, as the next result shows. In Section 3 we show how to use group representation theory to provide a proof of the matching upper bound.

Lemma 5. *For $c \geq 3$,*

$$\|\mathbb{P}(Y_{T_n-c} \in \cdot) - \pi_n(\cdot)\|_{\text{TV}} \geq 1 - 4^{1-c/3}.$$

Proof. In order to lower bound the total variation distance between Y and its equilibrium distribution at time $T_n - c$ we use the fact that the total variation distance is the maximal difference between the distribution of $Y_{T_n - c}$ and the uniform measure π_n on all possible subsets of \mathbb{Z}_n . So consider the set

$$A_n(c, \beta) = \{z \in \mathbb{Z}_n : |z - \mathbb{E}[Y_{T_n - c}]| > (3/8 - \beta)n\},$$

for some $\beta \in (0, 3/8)$ which we shall choose later. Note that this set satisfies $\pi_n(A_n(c, \beta)) = 1/4 + 2\beta$, and that (subject to this condition) it has been chosen to be as far away as possible from $\mathbb{E}[Y_{T_n - c}]$ (measured using the usual distance between two numbers mod n).

Using (2) we calculate the variance of Y'_k to be

$$\text{Var}[Y'_k] = \sum_{i=1}^k 4^{k+1-i} \sigma_{S'}^2 = \frac{4}{3}(4^k - 1)\sigma_{S'}^2, \quad (6)$$

and so

$$\begin{aligned} \mathbb{P}(Y_{T_n - c} \in A_n(c, \beta)) &\leq \mathbb{P}(|Y'_{T_n - c} - \mathbb{E}[Y'_{T_n - c}]| > (3/8 - \beta)n) \\ &\leq \frac{4^{1-c} n^2 \sigma_{S'}^2}{3\sigma_{S'}^2 ((3/8 - \beta)n)^2} = \frac{4^{1-c}}{3(3/8 - \beta)^2}. \end{aligned}$$

Here the first inequality follows from Y being equal to Y' mod n , and the second from Chebychev's inequality and the definition of T_n . Thus $A_n(c, \beta)$ satisfies

$$|\mathbb{P}(Y_{T_n - c} \in A_n(c, \beta)) - \pi_n(A_n(c, \beta))| \geq \frac{1}{4} + 2\beta - \frac{4^{1-c}}{3(3/8 - \beta)^2}. \quad (7)$$

This lower bound is maximised over values $\beta = \beta(c) \in (0, 3/8)$ when

$$\beta(c) = \frac{3}{8} - \left(\frac{4^{1-c}}{3}\right)^{1/3} \quad (c \geq 3)$$

and using this value of $\beta(c)$ yields the claimed left hand window of the cutoff:

$$\|\mathbb{P}(Y_{T_n - c} \in \cdot) - \pi_n(\cdot)\|_{\text{TV}} \geq 1 - \left(\frac{9}{4^{c-1}}\right)^{1/3} \geq 1 - 4^{1-c/3}, \quad c \geq 3.$$

□

3 Upper bound

3.1 Upper bounds and representation theory

Our basic method for obtaining upper bounds on the mixing times of our processes is to employ the techniques developed by [Diaconis and Shahshahani \(1981\)](#) for analysing random walks on groups. We briefly recall the main details. Given a finite group G , a (complex) representation ρ of G is a group homomorphism $\rho : G \rightarrow \text{GL}_n(\mathbb{C})$, where $\text{GL}_n(\mathbb{C})$ denotes the group of $n \times n$ invertible complex matrices. We call the number n the degree of ρ , denoted $\text{deg}(\rho)$, and we call the representation irreducible (or simple) if it cannot be decomposed into

a direct sum of two representations of smaller degree. Up to isomorphism, there are only finitely many such irreducible representations, and these include the trivial representation of degree 1 which sends every element of G to the complex number 1.

Given a probability P on G and a representation ρ , we can form the Fourier transform $\hat{P}(\rho)$ of P at ρ by setting

$$\hat{P}(\rho) := \sum_{g \in G} P(g)\rho(g),$$

so $\hat{P}(\rho)$ is an $n \times n$ matrix, where $n = \deg(\rho)$. One of the most attractive features of this Fourier transform is that it is well-behaved with respect to convolution, in that $\widehat{(P * Q)}(\rho) = \hat{P}(\rho)\hat{Q}(\rho)$ for any probabilities P and Q on G and any representation ρ . The following Upper Bound Lemma (Diaconis, 1988) allows one to compute an explicit upper bound for the total variation distance between a probability Q on G and the uniform distribution π on G . Since the Fourier transform behaves well with respect to convolution, this lemma provides a practical tool for bounding the mixing time of a random walk on a group.

Lemma 6. *Given a probability Q on a finite group G , we have*

$$\|Q - \pi\|_{\text{TV}}^2 \leq \frac{1}{4} \sum \deg(\rho) \text{tr}(\hat{Q}(\rho)\hat{Q}(\rho)^*),$$

where $A^* = (\overline{a_{ji}})$ denotes the complex conjugate transpose of the matrix $A = (a_{ij})$, tr denotes the trace function on square matrices, and the sum is taken over all non-trivial irreducible representations ρ of G .

3.2 Application to our walks

Our initial walk X involves both the additive and multiplicative structure of the ring \mathbb{Z}_n , and the measure giving the distribution of X_k cannot conveniently be expressed as the convolution of measures. This is the main reason we introduce the subsampled walk Y ; although Y is not strictly a random walk on the additive group $(\mathbb{Z}_n, +)$, the measure giving the distribution of Y_k can be expressed as the convolution of measures, and the techniques described in the previous section apply. Here the representation theory is particularly straightforward: there are precisely n irreducible representations $\rho_0, \rho_1, \dots, \rho_{n-1}$, they all have degree 1, and they are completely determined by the following equations

$$\rho_s(1) := e^{i\frac{2\pi}{n}s} \text{ for } 0 \leq s \leq n-1$$

(note that ρ_0 is the trivial representation). Therefore, for any probability Q on G and for any $0 \leq s \leq n-1$, $\hat{Q}(\rho_s)$ is just a complex number, $\hat{Q}(\rho_s)^*$ is just the complex conjugate of $\hat{Q}(\rho_s)$, and hence $\text{tr}(\hat{Q}(\rho_s)\hat{Q}(\rho_s)^*) = |\hat{Q}(\rho_s)|^2$. The Upper Bound Lemma becomes

$$\|Q - \pi\|_{\text{TV}}^2 \leq \frac{1}{4} \sum_{s=1}^{n-1} |\hat{Q}(\rho_s)|^2. \quad (8)$$

Recall from (2) that (with $Y_0 = 0$), $Y_k = \sum_{j=1}^k 2^j S'_{k+1-j} \pmod n$. The measure P_k giving the distribution of Y_k is the convolution of the measures λ_j given by $\lambda_j(2^j a \pmod n) = \mathbb{P}(S_1 = a)$ for every j, a , so we begin by calculating the Fourier transforms of the λ_j . To ease notation, for each $1 \leq j \leq k$ and $0 \leq s \leq n-1$, set

$$\omega_{sj} = \rho_s(2^j) = e^{i\frac{2\pi}{n}2^j s}$$

and note that for any j, s we have $\omega_{sj}^n = 1$. Then for each $0 \leq s \leq n-1$,

$$\begin{aligned}\hat{\lambda}_j(\rho_s) &= \sum_{a=0}^{n-1} \omega_{sj}^a \mathbb{P}(S_1 = a) = \sum_{a=0}^{n-1} \omega_{sj}^a \sum_{d \in \mathbb{Z}} \mathbb{P}(S'_1 = a + dn) \\ &= \sum_{d \in \mathbb{Z}} \sum_{a=0}^{n-1} \omega_{sj}^{a+dn} \mathbb{P}(S'_1 = a + dn) = \sum_{a \in \mathbb{Z}} \omega_{sj}^a \mathbb{P}(S'_1 = a) = G_{S'}(\omega_{sj}),\end{aligned}$$

where $G_{S'}$ is the probability generating function (PGF) of S' . It follows from its definition in (3) as a random sum of random step lengths that this satisfies

$$G_{S'}(\omega_{sj}) = \frac{p_n}{1 - (1 - p_n)G_\xi(\omega_{sj})}, \quad (9)$$

where G_ξ is the PGF of ξ .

When we substitute into the Upper Bound Lemma 6, we are interested in the modulus squared of such expressions, by Equation (8). The modulus of the top line squared is p_n^2 , and the modulus of the bottom line squared is

$$\begin{aligned}(1 - (1 - p_n)G_\xi(\omega_{sj}))(1 - (1 - p_n)\overline{G_\xi(\omega_{sj})}) \\ = 1 - (1 - p_n) \left(G_\xi(\omega_{sj}) + \overline{G_\xi(\omega_{sj})} \right) + (1 - p_n)^2 G_\xi(\omega_{sj}) \overline{G_\xi(\omega_{sj})} \\ = 1 - 2(1 - p_n) \operatorname{Re}(G_\xi(\omega_{sj})) + (1 - p_n)^2 |G_\xi(\omega_{sj})|^2.\end{aligned}$$

Combining all of the above leads to the following upper bound for the total variation distance at time k :

$$\|\mathbb{P}(Y_k \in \cdot) - \pi(\cdot)\|_{\text{TV}}^2 \leq \frac{1}{4} \sum_{s=1}^{n-1} \prod_{j=1}^k \frac{p_n^2}{1 - 2(1 - p_n) \operatorname{Re}(G_\xi(\omega_{sj})) + (1 - p_n)^2 |G_\xi(\omega_{sj})|^2}. \quad (10)$$

3.3 Strategy for analysing the upper bound

In order to establish a cutoff for Y , we need to control the right hand side of (10) around time $T_n = \log_2(n/\sigma_{S'})$. To that end, we define for $c \in \mathbb{N}$ a function $U_n(c)$ by

$$U_n(c) = \sum_{s=1}^{n-1} \prod_{j=1}^{T_n+c} \phi_n(s, j) \quad (11)$$

where

$$\phi_n(s, j) := \frac{p_n^2}{1 - 2(1 - p_n) \operatorname{Re}(G_\xi(\omega_{sj})) + (1 - p_n)^2 |G_\xi(\omega_{sj})|^2} \in (0, 1], \quad (12)$$

and note that our cutoff will be proved if we can show that (for odd n) $\limsup_{n \rightarrow \infty} U_n(c) \leq U(c)$ for some function U satisfying $U(c) \rightarrow 0$ as $c \rightarrow \infty$.

Our strategy for bounding $U_n(c)$ involves identifying for each $1 \leq s \leq n-1$ enough values j for which $\phi_n(s, j)$ is sufficiently small to provide a useful upper bound. In order to do this, it is convenient to first reparametrise, so we let Z_n be a random variable uniformly distributed on the set $\{s/n : s = 1, \dots, n-1\} \subset [0, 1]$. Then we may write

$$U_n(c) = \mathbb{E}[f_n(Z_n, T_n + c)], \quad \text{where } f_n(x, t) := (n-1) \prod_{j=1}^t \phi_n(nx, j). \quad (13)$$

The second step is to split the analysis of the function f_n into two stages by splitting the range of x into two pieces. In order to do this, let L be an integer satisfying $2\alpha L > 1$, and let b be an integer satisfying $B \subseteq [-2^b, 2^b]$, where (recall that) B is the support of ξ . We define a finite lattice \mathcal{L} of points in $[0, 1]$ by

$$\mathcal{L} = \left\{ \frac{k}{2^{L+b}} : k = 0, \dots, 2^{L+b} \right\}.$$

Now choose some $\varepsilon \in (0, 1/(2^{L+b}))$, and define the set \mathcal{L}_ε to be the intersection of $[0, 1]$ with

$$\bigcup_{x \in \mathcal{L}} \left[x - \frac{\varepsilon}{2}, x + \frac{\varepsilon}{2} \right].$$

Importantly, \mathcal{L}_ε depends only on α , B and ε , but not on n . We now proceed to bound $f_n(x, T_n + c)$ by considering in turn the cases where x does and does not belong to the set \mathcal{L}_ε .

3.4 Controlling f_n for $x \notin \mathcal{L}_\varepsilon$

For $x \notin \mathcal{L}_\varepsilon$ we see that $2\pi 2^j a x \not\equiv 0 \pmod{2\pi}$ for any $j = 1, 2, \dots, L$ and $a \in B$. Thus $\cos(2\pi 2^j a x)$ is bounded away from 1 for all such x and j , and we can write

$$\operatorname{Re} \left(G_\xi(e^{i2\pi 2^j x}) \right) = \sum_{a=0}^{2^b} \mathbb{P}(|\xi| = a) \cos(2\pi 2^j a x) \leq 1 - \kappa(x),$$

for all $j = 1, \dots, L$, where $\kappa(x)$ is strictly positive.

Substituting this into the expression for ϕ_n in (12), and lower-bounding the modulus squared of a complex number by the square of its real part, we obtain:

$$\begin{aligned} \phi_n(nx, j) &\leq \frac{p_n^2}{1 - 2(1 - p_n)\operatorname{Re}(G_\xi(e^{i2\pi 2^j x})) + (1 - p_n)^2 |G_\xi(e^{i2\pi 2^j x})|^2} \\ &\leq \frac{p_n^2}{1 - 2(1 - p_n)\operatorname{Re}(G_\xi(e^{i2\pi 2^j x})) + (1 - p_n)^2 \operatorname{Re}(G_\xi(e^{i2\pi 2^j x}))^2} \\ &= \left(\frac{p_n}{1 - (1 - p_n)\operatorname{Re}(G_\xi(e^{i2\pi 2^j x}))} \right)^2 \\ &\leq \left(\frac{p_n}{1 - (1 - p_n)(1 - \kappa(x))} \right)^2 = O(p_n^2). \end{aligned}$$

Since $\phi_n(nx, j) \in (0, 1]$, it follows that for $x \notin \mathcal{L}_\varepsilon$,

$$f_n(x, T_n + c) = (n-1) \prod_{j=1}^{T_n+c} \phi_n(nx, j) \leq (n-1) \prod_{j=1}^L \phi_n(nx, j) \leq O(n^{1-2\alpha L}).$$

Thanks to our choice of $L > 1/2\alpha$ we can now use Fatou's Lemma to deduce that

$$\limsup_{n \rightarrow \infty} \mathbb{E}[f_n(Z_n, T_n + c); Z_n \notin \mathcal{L}_\varepsilon] = 0. \quad (14)$$

3.5 Controlling f_n for $x \in \mathcal{L}_\varepsilon$

It remains to deal with $\mathbb{E}[f_n(Z_n, T_n + c); Z_n \in \mathcal{L}_\varepsilon]$. We begin by writing (for any $t \in \mathbb{N}$)

$$\begin{aligned} \mathbb{E}[f_n(Z_n, t); Z_n \in \mathcal{L}_\varepsilon] &= \frac{1}{n-1} \sum_{k=1}^{2^{L+b}-1} \sum_{r \geq 1} f_n\left(\frac{r}{n}, t\right) \mathbf{1}_{\left[\left|\frac{k}{2^{L+b}} - \frac{r}{n}\right| \leq \frac{\varepsilon}{2}\right]} \\ &\quad + \frac{1}{n-1} \sum_{r \geq 1} \left(f_n\left(\frac{r}{n}, t\right) + f_n\left(1 - \frac{r}{n}, t\right)\right) \mathbf{1}_{\left[\frac{r}{n} \leq \frac{\varepsilon}{2}\right]}, \end{aligned} \quad (15)$$

where the last sum deals with the two end intervals in \mathcal{L}_ε . Now (15) can be bounded as follows:

$$\mathbb{E}[f_n(Z_n, t); Z_n \in \mathcal{L}_\varepsilon] \leq \frac{1}{n-1} \sum_{k=1}^{2^{L+b}-1} \sum_{r=-\infty}^{\infty} f_n\left(\frac{(n-1)\frac{k}{2^{L+b}} - r}{n}, t\right) + \frac{2}{n-1} \sum_{r=1}^{\infty} f_n\left(\frac{r}{n}, t\right), \quad (16)$$

where we have used the symmetry of the functions f_n at either end of the interval $[0, 1]$ to rewrite the expression for the end intervals. Now replace t by $T_n + c$, and consider the function f_n in the double sum above:

$$\begin{aligned} f_n\left(\frac{(n-1)\frac{k}{2^{L+b}} - r}{n}, T_n + c\right) &= (n-1) \prod_{j=1}^{T_n+c} \phi_n\left((n-1)\frac{k}{2^{L+b}} - r, j\right) \\ &\leq (n-1) \phi_n\left((n-1)\frac{k}{2^{L+b}} - r, T_n + c\right). \end{aligned} \quad (17)$$

Here we have bounded the product by a single term, once again making use of the fact that ϕ_n takes values in $(0, 1]$. Since $\phi_n(s, j)$ involves s and j only through the function $G_\xi(\omega_{sj})$, where $\omega_{sj} = \exp(2\pi i 2^j s/n)$, we have (for sufficiently large n) that the bound in (17) is a function of

$$\begin{aligned} \exp\left(2\pi i \frac{2^{T_n+c}}{n} \left((n-1)\frac{k}{2^{L+b}} - r\right)\right) &= \exp\left(2\pi i \frac{2^{T_n+c}}{n} (k2^{-(L+b)} + r)\right) \\ &= \exp\left(\frac{2^{1+c}\pi i (k2^{-(L+b)} + r)}{\sigma_{S'}}\right). \end{aligned}$$

The second equality simply uses the definition of T_n , while the first results from shifting the argument of the exponential function by $2\pi i k 2^{T_n+c-(L+b)}$. (For large enough n this is an integer multiple of $2\pi i$, thanks to the finiteness of L and b and the assumption that $T_n \rightarrow \infty$.)

Writing

$$\theta_{krc} = \frac{2^{1+c}\pi (k2^{-(L+b)} + r)}{\sigma_{S'}},$$

we therefore need to upper bound the function

$$\phi_n\left((n-1)\frac{k}{2^{L+b}} - r, T_n + c\right) = \frac{p_n^2}{1 - 2(1-p_n)\operatorname{Re}(G_\xi(e^{i\theta_{krc}})) + (1-p_n)^2 |G_\xi(e^{i\theta_{krc}})|^2}.$$

Now note that

$$G_\xi\left(e^{i\theta_{krc}}\right) = \sum_{a \in B} \mathbb{P}(\xi = a) e^{ia\theta_{krc}}, \quad \text{and thus} \quad \operatorname{Re}\left(G_\xi\left(e^{i\theta_{krc}}\right)\right) = \mathbb{E}[\cos(\xi\theta_{krc})].$$

Similarly,

$$\left| G_\xi(e^{i\theta_{krc}}) \right|^2 = \mathbb{E} [\cos(\xi\theta_{krc})]^2 + \mathbb{E} [\sin(\xi\theta_{krc})]^2 .$$

Since $p_n \rightarrow 0$ as $n \rightarrow \infty$, we see from (5) that $\sigma_{S'} \rightarrow \infty$ and thus $\theta_{krc} \rightarrow 0$. Using the Taylor expansions of cosine and sine the above can be approximated by

$$\begin{aligned} \mathbb{E} [\cos(\xi\theta_{krc})] &= 1 - \frac{(\mu^2 + \sigma^2)\theta_{krc}^2}{2} + O(\theta_{krc}^4); \\ \mathbb{E} [\sin(\xi\theta_{krc})] &= \mu\theta_{krc} + O(\theta_{krc}^3). \end{aligned}$$

Neglecting terms of $O(\theta_{krc}^3)$ we arrive at

$$\begin{aligned} \phi_n \left((n-1)\frac{k}{2^{L+b}} - r, T_n + c \right) &\sim \frac{p_n^2}{1 - (1-p_n) \left[2 - (\mu^2 + \sigma^2)\theta_{krc}^2 \right] + (1-p_n)^2 \left[1 - (\mu^2 + \sigma^2)\theta_{krc}^2 + \mu^2\theta_{krc}^2 \right]} \\ &= \frac{p_n^2}{p_n^2 + (1-p_n)(\mu^2 + \sigma^2 p_n)\theta_{krc}^2} \\ &= \frac{1}{1 + \sigma_{S'}^2 \theta_{krc}^2} = \frac{1}{1 + 4^{1+c} \pi^2 (k2^{-(L+b)} + r)^2} . \end{aligned}$$

We now combine this bound with that in (17) and insert into (16) (using an identical argument for the second sum there):

$$\mathbb{E} [f_n(Z_n, t); Z_n \in \mathcal{L}_\varepsilon] \leq \sum_{k=1}^{2^{L+b}-1} \sum_{r=-\infty}^{\infty} \frac{1}{1 + 4^{1+c} \pi^2 (k2^{-(L+b)} + r)^2} + 2 \sum_{r=1}^{\infty} \frac{1}{1 + 4^{1+c} \pi^2 r^2} .$$

Summing over r this bound becomes

$$\mathbb{E} [f_n(Z_n, T_n + c); Z_n \in \mathcal{L}_\varepsilon] \leq \sum_{k=1}^{2^{L+b}-1} \frac{2^{-(1+c)} \sinh(2^{-c})}{\cosh(2^{-c}) - \cos(2^{1-(L+b)} k \pi)} + \left(2^{-(1+c)} \coth(2^{-(1+c)}) - 1 \right) .$$

Finally, taking the worst possible value $k = 1$ for *every* term in the sum, we get

$$\begin{aligned} \mathbb{E} [f_n(Z_n, T_n + c); Z_n \in \mathcal{L}_\varepsilon] &\leq \frac{2^{L+b-c} \sinh(2^{-c})}{\cosh(2^{-c}) - \cos(2^{1-(L+b)} \pi)} + \left(2^{-(1+c)} \coth(2^{-(1+c)}) - 1 \right) \\ &\sim \frac{2^{L+b-2c}}{1 - \cos(2^{1-(L+b)} \pi)} + \frac{4^{1-c}}{3} \quad \text{as } c \rightarrow \infty. \end{aligned} \tag{18}$$

Combining (14) and (18) yields the required result

$$\limsup_{n \rightarrow \infty} U_n(c) = \limsup_{n \rightarrow \infty} \mathbb{E} [f_n(Z_n, T_n + c); Z_n \in \mathcal{L}_\varepsilon] \leq O(4^{-c}) \quad \text{as } c \rightarrow \infty,$$

and this completes the proof of Theorem 3.

4 Concluding remarks

We have shown that the process X exhibits a cutoff phenomenon when the probability p_n of jumping takes the form $p_n = 1/(2n^\alpha)$, for a range of α which depends upon the mean of

our step distribution ξ ($\alpha \in (0, 2)$ when $\mu = 0$, and $\alpha \in (0, 1)$ otherwise). We have not yet said anything about the mixing time when α takes values on the boundary of these intervals, however.

If $\alpha = 0$ then our argument for upper bounding the mixing time breaks down (since a sufficiently fine lattice \mathcal{L} does not exist). In this situation Lemma 5 is still applicable, showing that Y has not mixed by time $\log_2 n$, and an upper bound of $O(\log_2 n \log_2 \log_2 n)$ can be obtained by employing the method of Chung et al. (1987). On the other hand, if α takes the value at the upper boundary of the relevant interval then $n/\sigma_{S'} = O(1)$, and thus T_n is asymptotically independent of n : in this case our upper bound analysis still holds, and we see that Y mixes in constant time.

It is of course possible to generalise the process considered in this paper in a number of ways. Changing the form of p_n to $1/(cn^\alpha)$ for some constant $c > 1$ has no significant effect on the mixing time of Y , and so X will exhibit a cutoff at time $cn^\alpha \log_2(n/\sigma_{S'})$. Similarly, changing the transitions of X so that jumps involve multiplying by some (fixed) $k \geq 2$ (and considering only those n for which the resulting process still has a uniform equilibrium distribution) presumably just has the effect of changing the cutoff time for Y to $\log_k(n/\sigma_{S'})$. More interesting would be an analysis of a process X for which the multiplication factor is not deterministic, and for which the resulting subsampled chain Y does not have a distribution given by simple convolution; for example where $X_t = a_t X_{t-1}$ with probability p_n , with a_t being uniformly chosen from $\{2, (n+1)/2\}$.

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Department of Mathematics
University of York
York, YO10 5DD
UK