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On the Outage Performance of SWIPT Based Three-step Two-way DF Relay Networks

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Abstract—In this paper, we study the outage performance of simultaneous wireless information and power transfer (SWIPT) based three-step two-way decode-and-forward (DF) relay networks, where both power-splitting (PS) and “harvest-then-forward” are employed. In particular, we derive the expressions of terminal-to-terminal (T2T) and system outage probabilities based on a Gaussian-Chebyshev quadrature approximation, and obtain the T2T and system outage capacities. The effects of various system parameters, e.g., the static power allocation ratio at the relay, symmetric PS, as well as asymmetric PS, on the outage performance of the investigated network are examined. It is shown that our derived expression for T2T outage capacity is more accurate than existing analytical results, and that the asymmetric PS achieves a higher system outage capacity than the symmetric one when the channels between the relay node and the terminal nodes have different statistic gains.

Index Terms—Outage performance, SWIPT based three-step two-way DF relay, power-splitting.

I. INTRODUCTION

Simultaneous wireless information and power transfer (SWIPT) has been recognized as a promising technology to alleviate the energy constraint in relay networks while maintaining reliable communications through a power-splitting (PS) or time-switching (TS) scheme. In a bidirectional transmission network, two-way relaying operates in two steps or three steps instead of four steps as required in one-way relaying, thus achieving a higher capacity. In this regard, SWIPT based three-step (or two-step) two-way relay networks (TWRNs) are receiving increasing attention from the wireless industry and academia. Due to the limited space and our attention on SWIPT based three-step TWRNs, our work omits the literature review of SWIPT based two-step TWRNs, which can be referred to as the references.

As the low complexity of hardware is very important to energy-constrained networks, while the circuit of three-step TWRNs is simpler than that of two-step TWRNs, three-step TWRNs have attracted vast attention recently. The authors of derived the terminal-to-terminal (T2T) outage probability, which considers the outage events at the terminal nodes independently, for three wireless power transfer schemes in TS based three-step amplify-and-forward (AF) TWRNs. Considering three-step AF multiplicative TWRNs, the authors in proposed a symmetric PS scheme and investigated the T2T outage probability. To minimize the system outage probability, an optimal symmetric PS scheme and an optimal asymmetric scheme were developed in and , respectively. Recently, three-step two-way relaying has been extended to DF relaying, where both upper and lower bounds of the T2T outage capacity were obtained for the symmetric PS scheme. The authors of derived expressions of system outage probability for three-step TWRNs with TS(PS) SWIPT. To explore the advantages of both AF and DF relaying protocols, the authors studied the T2T performance of a SWIPT based three-step TWR with hybrid decode-amplify-forward (HDAF). The combination of SWIPT based three-step with other cognitive radio was reported in , aiming to analyze the system outage probability for the primary user and the secondary user.

Although these works mentioned above have laid a solid foundation for the understanding of SWIPT based three-step TWRNs, there are still large gaps needed to be filled. In particular, for a SWIPT based three-step DF TWRN, the existing theoretical expression of T2T outage capacity is far from the simulation results. Besides, the authors in and focused on the analysis of system outage probability, while they ignored the channel reciprocity of terminal-relay links when deriving the expression of system outage probability. Accordingly, the system outage performance of SWIPT based three-step DF TWRNs is still far from being well understood. Motivated by the above observations, in this paper, we study the T2T and system outage probabilities of three-step DF TWRNs, where PS and “harvest-then-forward” are employed at the relay node, and derive expressions for both of them. Different from the T2T outage probability, the system outage probability evaluates the overall outage performance of both T2T links (i.e., the two one-way relay links) in the

1In PS based three-step TWRNs, the relay is equipped with one power splitter and the PS ratios for both terminal-relay links may be unequal. If the PS ratios for two terminal-relay links are always equal, we refer to the PS scheme as the symmetric PS; otherwise, we refer to it as the asymmetric PS.
TWRNs, making the analysis and derivation much challenging than those of T2T outage performance. Based on the derived expressions, the impacts of various system parameters, such as the relay location and the energy conversion efficiency, on the system outage performance are analyzed. The performance comparison between symmetric and asymmetric PS schemes is also presented.

II. SYSTEM MODEL

We consider a DF TWRN, where two terminal nodes A and B exchange information with the aid of an energy-constrained relay node R based on a PS scheme, as shown in Fig. 1. The whole block time $T$ has three transmission slots, where $\beta T$ ($0 < \beta < 0.5$) is the time proportion for the relay to harvest energy and decode signal from one terminal node [10]. The “harvest-then-forward” scheme is adopted to encourage the relay to assist two terminal nodes. All nodes are equipped with a single antenna and operate in a half-duplex mode. We assume that there is no direct link between A and B. Note that our considered network can be applied to many energy-constrained wireless scenarios, e.g., the wireless sensor network in a toxic environment and the wireless body area network [11].

The channel model is given by $|h_i|^2d_\alpha (i = A \text{ or } B)$, where $h_i \sim \mathcal{CN}(0, \mu_i)$ is the $i$–R channel fading coefficient; $d_i$ is the $i$–R distance; $\alpha$ is the path loss exponent. Thus the statistic channel gain of $i$–R link is given by $\mu_i d_\alpha$. For analytical tractability, we further ignore the consumption of the transceiver circuitry at R [10].

At the first and the second time slots (each of the time slot is $\beta T$), nodes A and B transmit the normalized information signals $x_A$ and $x_B$ to the relay, respectively. The received signal at the relay R from terminal node i is $y_{i\rightarrow R} = h_i \sqrt{P_i} d^\alpha (i = A \text{ or } B)$, where $n_{Ri} \sim \mathcal{CN}(0, \sigma^2_{Ri})$ is the additive white Gaussian noise (AWGN) and $P_i$ denotes the transmit power of terminal node i. Here, we assume that $P_A = P_B = P_0$. Using the PS scheme, the received signal $y_{i\rightarrow R}$ is split to two parts: $\sqrt{\lambda_i}y_{i\rightarrow R}$ for energy harvesting EH) and $\sqrt{1-\lambda_i}y_{i\rightarrow R}$ for information processing. Thus, the total received energy from both terminal nodes and the received SNR from i to R are given by $E_{\text{total}} = P_0 \eta \beta T (\lambda_i \|h_i\|^2 d^{-\alpha}_A + \lambda_B \|h_B\|^2 d^{-\alpha}_B)$ and $\gamma_{i\rightarrow R} = P_0 \beta T (1 - \lambda_i) d^{-\alpha}_A \eta^{-2}$, respectively, where $\eta$ is the energy conversion efficiency and $\lambda_i$ is the PS ratio of terminal node i. Note that $\lambda_B$ can be different from $\lambda_A$ in three-step relay networks, which makes it possible to more efficiently use of the received radio frequency (RF) signal at the relay.

Let $\bar{x}_i$ be the decoded signal for $x_i$. If both $x_A$ and $x_B$ are successfully decoded during the first two slots, then at the third time slot $(1 - 2\beta)T$ the relay R broadcasts the normalized signal $x_R = \theta_A x_A + \theta_B x_B$ to both terminal nodes using the harvested energy $E_{\text{total}}$, where $\theta^2_A + \theta^2_B = 1$ and $\theta^2_i (i = A, B)$ is a static power allocation ratio determining how the relay allocates the power for the decoded signals based on the statistic channel state information (CSI) (or statistic channel gain) of the i–R link. Thus, the received signal at terminal node i is given by $y_{R\rightarrow i} = h_i \sqrt{P_R d^\alpha} x_R + n_{Ri} \sim \mathcal{CN}(0, \sigma^2_{Ri})$, where if $i = A, \bar{i} = B$; $i = B, \bar{i} = A$; $n_{Ri} \sim \mathcal{CN}(0, \sigma^2_{Ri})$ is the AWGN; $P_R = \frac{\rho_0}{\beta \sigma^2}$ is the transmit power at R; step (a) holds under the assumption of perfect successive interference cancellation (SIC) [10]–[12]. This assumption is used here in order to obtain the upper bound of outage performance for our considered network.

Accordingly, the SNR of the $R \rightarrow i$ link is given by $\gamma_{R\rightarrow i} = \frac{\sigma^2_{Ri}(\lambda_A \|h_A\|^2 d^{-\alpha}_A + \lambda_B \|h_B\|^2 d^{-\alpha}_B)}{\beta \sigma^2}$, where $\rho_0 = \frac{\rho_0}{\beta \sigma^2}$ and $\sigma^2_{Ri} = \sigma^2_{RA} + \sigma^2_{RB} = \sigma^2_{RA} = \sigma^2_{RB} = \sigma^2$.

III. OUTAGE PROBABILITY ANALYSIS

A. Terminal-to-Terminal Outage Probability

Let $P(\cdot)$ denote the probability and $P_{out}$ be the outage probability of $i \rightarrow R \rightarrow i$ link. For a given target transmission rate $U$ and the corresponding SNR threshold $\gamma_{th}$ [10], $P_{out}$ is written as

$$P_{out} = 1 - P_i$$ (1)

where $P_i = P(\gamma_{R\rightarrow i} \geq \gamma_{th}, \gamma_{i\rightarrow R} \geq \gamma_{th})$.

Based on the expression of $\gamma_R \rightarrow i$ and $\gamma_{i\rightarrow R}$, $P_i$ can be rewritten as

$$P_i = P(\|h_i\|^2 \geq \gamma_{th} (\Phi_i, \Psi_i)) = P(\|h_i\|^2 \geq \Psi_i, \Phi_i \geq \Psi_i) + P(\|h_i\|^2 \geq \Phi_i, \Phi_i > \Psi_i)$$ (2)

When $\Phi_i > \Psi_i$, we have the following inequality

$$a_i |h_i|^4 + b_i |h_i|^2 - c_i \leq 0$$ (3)

where $a_i = \lambda_i d^{-\alpha}_A, b_i = \frac{\lambda_B}{\mu_i(1-\lambda_i)}$, and $c_i = \gamma_{th}/X_i$.

Combining (3) with $|h_i|^2 \geq 0$, the range of $|h_i|^2$ is $0 \leq |h_i|^2 \leq \Omega_i$, where $\Omega_i = \sqrt{\beta \sigma^2 + 4a_i c_i} / b_i$.

When $\Phi_i > \Psi_i$, we have $|h_i|^2 > \Omega_i$. Thus, $P_i$ can be calculated as

$$P_i = P(\|h_i\|^2 > \Phi_i, |h_i|^2 > \Omega_i) + P(\|h_i\|^2 > \Psi_i, 0 \leq |h_i|^2 \leq \Omega_i)$$

$$= \exp \left( -\Phi_i \frac{\Omega_i^2}{\mu_i} \right) + \frac{1}{\mu_i} \int_0^{\Omega_i} \exp \left( -\Psi_i y + y \frac{y}{\mu_i} \right) dy$$ (4)

where the second equality holds from $|h_i|^2 \sim \exp(\frac{1}{\mu_i})$ for $i \in \{A, B\}$ and $y = |h_i|^2$.

Since it is difficult to find the closed-form expression for $P_i$ due to the integral $\int_0^{\Omega_i} \exp(z_1 x + z_2) dx$ for any value of

$2$The assumption will not cause any loss of generality to our analysis. This is because the average signal-to-noise ratios (SNRs) of all channels are essential for the analysis of system outage performance, and the average SNRs depend on the transmit power at terminals and all the fading channels. Although $P_A = P_B$, the independent fading channels will have different values for their parameters, e.g., the distance of the terminal-relay link, making the average SNRs of all channels different.

3Note that the system capacity can be further improved if a dynamic power allocation ratio, which is adjusted based on the instantaneous CSI (or instantaneous channel gain) of terminal-relay links instead of static one, is adopted. This will be considered in our future work.
where $N$ is a parameter that determines the tradeoff between complexity and accuracy for the Gaussian-Chebyshev quadrature based approximation; $\nu_n = \cos \frac{2n-1}{2N}\pi$, and $\chi_n^{(1)} = \frac{\nu_n}{\mu_n} + \frac{\nu_0}{2}$. Note that an acceptable accuracy can be achieved for a small value of $N$, which is verified in our simulations.

Substituting (5) into (1), we obtain the T2T outage probability for $\bar{i} \rightarrow R \rightarrow i$ link, $P_{i}^{1\bar{i}}$.

Thus, the outage capacity $\tau_{out}^{i}$ can be calculated as

$$\tau_{out}^{i} = (1 - P_{out}^{i}) U \beta T. \tag{6}$$

### B. System Outage Probability

The system outage probability is defined as the probability that either the SNR of $i \rightarrow i$ link is less than the SNR threshold $\gamma_{th}$. Thus, the system outage probability is written as

$$P_{out}^{S} = 1 - \mathbb{P}(\cup_{i=1:K} \gamma_i \geq \gamma_{th}) \tag{7}$$

where $P_{i}^{S} = 1 - \mathbb{P}(\cup_{i=1:K} \gamma_i \geq \gamma_{th})$. Based on the expressions of $\psi_i$ and $\Psi_i$ defined in (2), $P_{i}^{S}$ is rewritten as

$$P_{i}^{S} = P_{i1}^{S} + P_{i2}^{S} + P_{i3}^{S} + P_{i4}^{S} \tag{8}$$

where

$$P_{i1}^{S} = \mathbb{P}(|h_i|^2 \geq \psi_i, |h_i|^2 \geq \psi_B, \Psi_A \geq \Psi_B, \Psi_B < \Phi_B) \tag{9}$$

$$P_{i2}^{S} = \mathbb{P}(|h_i|^2 \geq \psi_B, |h_i|^2 \geq \psi_A, \Psi_A < \Psi_B, \Psi_B < \Phi_B) \tag{10}$$

$$P_{i3}^{S} = \mathbb{P}(|h_i|^2 \geq \psi_A, |h_i|^2 \geq \psi_A, \Psi_A < \Psi_B, \Psi_B < \Phi_B) \tag{11}$$

$$P_{i4}^{S} = \mathbb{P}(|h_i|^2 \geq \psi_A, |h_i|^2 \geq \psi_A, \Psi_A < \Psi_B, \Psi_B < \Phi_B). \tag{12}$$

1) $P_{i1}^{S}$: If $\Psi_A \geq \Psi_B$ and $\Psi_B < \Phi_B$, we have

$$P_{i1}^{S} = \mathbb{P}(|h_i|^2 \geq \psi_A, |h_i|^2 \geq \psi_B, \Psi_A \geq \Psi_B, \Psi_B < \Phi_B) \tag{13}$$

Since $|h_i|^2 \geq \psi_A$ and $|h_i|^2 \geq \psi_B$, the solution to (9) is given by

$$P_{i1}^{S} = \frac{b_i |h_i|^2}{2a_i} \left( |h_i|^2 \leq \Omega_B \right)$$

where $a_i = |h_i|^2 - c_i > 0$.

2) $P_{i2}^{S}$: If $\Psi_A < \Phi_B$ and $\Psi_B \geq \Phi_B$, we have $0 \leq |h_A|^2 \leq \Omega_A$ and $|h_B|^2 > \Omega_B$. Similar to $P_{i1}^{S}$, $P_{i2}^{S}$ can be calculated as

$$P_{i2}^{S} \approx \frac{\pi \max(\Phi_A, \Omega_A) - \Phi_A}{2N\mu_B} \left( \sum_{n=1}^{N} \frac{1 - \nu_n^2}{\mu_n} \right) \tag{14}$$

where $\phi_B(x) = \max(\Psi_B, \Omega_B)$ with $|h_A|^2 = x$ and $\chi_n^{(2)} = (\max(\Phi_A, \Omega_A) - \Phi_A) \nu_n + (\max(\Phi_A, \Omega_A) + \Phi_A)\mu_n$.

3) $P_{i3}^{S}$: If $\Psi_A < \Phi_B$ and $\Psi_B < \Phi_B$, the ranges of $|h_B|^2$ and $|h_A|^2$ are determined by $|h_B|^2 > \Omega_B$ and $|h_A|^2 > \Omega_A$, respectively. Combining these ranges with $|h_B|^2 \geq \Phi_B$ and $|h_A|^2 \leq \Phi_A$, $P_{i3}^{S}$ can be computed as

$$P_{i3}^{S} = \mathbb{P}(|h_B|^2 \geq \max(\Phi_B, \Omega_B), |h_A|^2 \geq \max(\Phi_A, \Omega_A)) \tag{15}$$

4) $P_{i4}^{S}$: If $\Psi_A \geq \Phi_A$ and $\Psi_B \geq \Phi_B$, both $0 \leq |h_B|^2 \leq \Omega_B$ and $0 \leq |h_A|^2 \leq \Omega_A$ should be satisfied. Then $P_{i4}^{S}$ can be denoted as

$$P_{i4}^{S} = \mathbb{P}(\Psi_A \geq |h_A|^2, \Psi_B \geq |h_B|^2) \tag{16}$$

We find that the integral region for calculating $P_{i4}^{S}$ is bounded by the following 4 lines: $y = \frac{\gamma_B}{X_{\lambda_{A_{xy}}}} - 2a_B^{\lambda_{A_{xy}}}$, $y = x$, $y = \frac{\gamma_A}{X_{\lambda_{A_{xy}}}} - 2a_B^{\lambda_{A_{xy}}}$, and $y = \frac{\gamma_B}{X_{\lambda_{A_{xy}}}} - 2a_B^{\lambda_{A_{xy}}}$.

Let $(x_1, y_1)$ denote the intersection of the lines $x = x_1$, $y = y_1$, and $x = \frac{\gamma_B}{X_{\lambda_{A_{xy}}}} - 2a_B^{\lambda_{A_{xy}}}$. Then $(x_1, y_1)$ is given by

$$x_1 = \frac{\frac{\gamma_A}{X_{\lambda_{A_{xy}}}} - \frac{\gamma_B}{X_{\lambda_{A_{xy}}}} + 4a_Bc_Bd_B \lambda_{A_{xy}}}{2a_B^{\lambda_{A_{xy}}}}, y_1 = \frac{\gamma_A}{X_{\lambda_{A_{xy}}}}. \tag{17}$$

Similarly, the intersection of the lines $y = \gamma_B$ and $y = \frac{\gamma_B}{X_{\lambda_{A_{xy}}}} - 2a_B^{\lambda_{A_{xy}}}$ is given by

$$x_2 = \frac{\frac{\gamma_B}{X_{\lambda_{A_{xy}}}} - \frac{\gamma_B}{X_{\lambda_{A_{xy}}}} + 4a_Bc_Bd_B \lambda_{A_{xy}}}{2a_B^{\lambda_{A_{xy}}}}, y_2 = \frac{\gamma_A}{X_{\lambda_{A_{xy}}}}. \tag{18}$$

The intersection of the lines $x = \gamma_B$ and $x = \frac{\gamma_B}{X_{\lambda_{A_{xy}}}} - 2a_B^{\lambda_{A_{xy}}}$ is denoted by $(q_1, y_1)$, where $q_1 = \frac{\gamma_B}{X_{\lambda_{A_{xy}}}} - 2a_B^{\lambda_{A_{xy}}}$.

The intersection of the lines $x = \frac{\gamma_B}{X_{\lambda_{A_{xy}}}} - 2a_B^{\lambda_{A_{xy}}}$ and $y = \frac{\gamma_B}{X_{\lambda_{A_{xy}}}} - 2a_B^{\lambda_{A_{xy}}}$ is denoted by $(x_1, q_2)$, where $q_2 = \frac{\gamma_B}{X_{\lambda_{A_{xy}}}} - 2a_B^{\lambda_{A_{xy}}}$.

The intersection of the lines $x = \frac{\gamma_B}{X_{\lambda_{A_{xy}}}} - 2a_B^{\lambda_{A_{xy}}}$ and $y = \frac{\gamma_B}{X_{\lambda_{A_{xy}}}} - 2a_B^{\lambda_{A_{xy}}}$ is determined by

$$x = C_A/y - D_A B y$$

$$y = C_B/x - D_B x$$

where $C_i = \gamma_A d_i/X_{\lambda_{A_{xy}}}$ and $D_i = \gamma_A d_i/X_{\lambda_{A_{xy}}}$. Substituting the $y = C_A/y - D_A B y$ into $x = C_A/y - D_A B y$, we have

$$- C_A + C_B)^2 + C_B^2 D_A = 0. \tag{19}$$

Since $x > 0$, the solution to (19) is given by $x_y = \sqrt{C_B^2 D_A/(C_A + C_B)}$, and the corresponding value of $y$ is $y_B = C_B/x - D_B x_0$. 

\footnote{Note that, compared with traditional three-step TWFRNs without SWIPT, the derivation of system outage probability in SWIPT based three-step TWFRNs is more challenging since the transmit power of the relay depends on all the channel gains of terminal-relay links.}
According to the values of intersections, the integral region for $P_{14}$ can be divided into three cases.

**Case I:** If $\max(q_1, x_\Delta) \geq x_1$ or $\max(q_2, y_\Delta) \geq y_1$, the integral region is $0 = 0$ and $P_{14} = 0$.

**Case II:** If $\max(q_1, x_\Delta) < x_1$, $\max(q_2, y_\Delta) < y_1$, and $x_0 \leq \max(q_1, x_\Delta)$ (or $x_0 > x_1$), the integral region for $P_{14}$ is bounded by the following three lines: $x = \Omega_1$, $y = \Omega_B$, and $y = C_B/x - D_Bx$ (or $x = C_A/y - D_Ay$).

When $y_\Delta \geq q_2$, the integral region for $P_{14}$ is shown in Fig. 2. Accordingly, $P_{14}$ is given by

$$P_{14}^s = \frac{1}{\mu_B} \int_{x_1}^{y_1} \exp \left( -\frac{C_A}{\mu_A} + \frac{D_A}{\mu_A} - \frac{1}{\mu_B} \right) y \ dy - \frac{1}{\mu_B} \int_{x_1}^{y_1} \exp \left( \frac{C_B}{\mu_B} + \frac{D_B}{\mu_B} - \frac{1}{\mu_B} \right) y \ dy$$

**Case III:** If $\max(q_1, x_\Delta) = 1$ and $\max(q_2, y_\Delta) = 1$, the integral region for $P_{14}$ is bounded by four lines.

For the case with $y_\Delta < q_2$, $P_{14}$ is given by

$$P_{14}^s = \frac{1}{\mu_A} \int_{x_1}^{y_1} \exp \left( \frac{C_B}{\mu_B} + \frac{D_B}{\mu_B} - \frac{1}{\mu_B} \right) y \ dy - \frac{1}{\mu_B} \int_{x_1}^{y_1} \exp \left( \frac{C_A}{\mu_A} + \frac{D_A}{\mu_A} - \frac{1}{\mu_B} \right) y \ dy - \frac{1}{\mu_B} \int_{x_1}^{y_1} \exp \left( \frac{C_B}{\mu_B} + \frac{D_B}{\mu_B} - \frac{1}{\mu_B} \right) y \ dy$$

Similarly, for the case with $y_\Delta < q_2$, $P_{14}$ is given by

$$P_{14} = \frac{1}{\mu_B} \int_{x_1}^{y_1} \exp \left( \frac{C_B}{\mu_B} + \frac{D_B}{\mu_B} - \frac{1}{\mu_B} \right) y \ dy - \frac{1}{\mu_B} \int_{x_1}^{y_1} \exp \left( \frac{C_A}{\mu_A} + \frac{D_A}{\mu_A} - \frac{1}{\mu_B} \right) y \ dy$$

where

$$\chi_n^{(5)} = \frac{(y_1 - y_\Delta)n + (y_1 + y_\Delta)}{2}$$

and

$$\chi_n^{(6)} = \frac{(y_2 - y_\Delta)n + (y_2 + y_\Delta)}{2}$$

Similarly, the system outage capacity $\tau_{out}$ is determined.

$$\tau_{out} = \frac{1 - P_{out}}{\mu_B}$$

**Remark:** The derived results can serve the following purposes. The first purpose is to obtain an accurate outage capacity with a small $N$ instead of the computer simulations. The second purpose is to observe some insights regarding the design of the PS ratios and the static power allocation ratio from the curves obtained by the derived results. In our considered network, the PS ratios and the static power allocation ratio are designed based on the statistic channel gains instead of instantaneous channel gains, it is practical to obtain the insights offline regarding the design of PS ratios and the static power allocation ratio and such approach has also adopted in many works, e.g., [6]. A concrete example is Fig. 5 of the revised manuscript, where the optimal system outage capacity and PS ratios for two PS schemes are obtained by maximizing the derived system outage capacity expression as in problem (23). It can be observed that the optimal $\lambda_i$ decreases with the increase of the statistic channel gain of $i-\text{R link}$.

$$\max_{\lambda_A, \lambda_B} \tau_{out} \ s.t. \ 0 \leq \lambda_A, \lambda_B \leq 1$$

Lastly, we can obtain the diversity gain of the considered network, shown in the next subsection.
Thus, the diversity gain can be calculated as
\[ d = \lim_{\rho_0 \to \infty} \frac{\log (1 - P_{11}^{s} - P_{12}^{s} - P_{13}^{s} - P_{14}^{s})}{\log (\rho_0)}. \]  
(24)

It is not hard to verify that \( \lim_{\rho_0 \to \infty} P_{11}^{s} = \lim_{\rho_0 \to \infty} P_{12}^{s} = \lim_{\rho_0 \to \infty} P_{14}^{s} = 0 \) and \( \lim_{\rho_0 \to \infty} P_{13}^{s} = 1 \). Thus, the diversity gain can be rewritten as
\[ d = - \lim_{\rho_0 \to \infty} \frac{\log (1 - P_{13}^{s})}{\log (\rho_0)}. \]  
(25)

Now let us compare \( \Omega_i \) and \( \Phi_i \). When the transmit SNR approaches infinite, we have
\[ \lim_{\rho_0 \to \infty} \frac{\Omega_i}{\Phi_i} = \infty. \]  
(26)

Combining (26) and the fact that \( \Omega_i \) and \( \Phi_i \) decrease with \( \rho_0 \), we have
\[ \lim_{\rho_0 \to \infty} \max(\Omega_i, \Phi_i) = \lim_{\rho_0 \to \infty} \Phi_i. \]  
(27)

Thus, the diversity gain can be calculated as
\[ d = - \lim_{\rho_0 \to \infty} \frac{\log (1 - \exp (-\frac{1}{\rho_0}))}{\log (\rho_0)}, \]  
\[ x = \frac{1}{\rho_0}, \quad \lim_{x \to 0} \frac{x}{1 - \exp (-x)} = 1. \]  
(28)

### IV. Simulations

During our simulation results, we provide an important validation of the derived expressions of (9) and (22). We adopt the same system parameter settings as in [10] for a fair comparison between [6] and the existing T2T outage capacity [10]. Fig. 4 plots the relative approximate error versus \( N \), as well as the outage capacity as a function of the PS ratio by assuming that \( \lambda_A = \lambda_B \) as in [10]. For the upper part of Fig. 4, we assume that \( \lambda_A = \lambda_B = 0.2 \). It can be seen that the relative approximation error [15] decreases with \( N \), and our approximated results are sufficient to provide level of accuracy within very few items (e.g. \( N \geq 5 \)). For the lower part of Fig. 4, the upper and lower bounds of the T2T outage probability from [10] are also provided. The theoretical T2T outage capacity and system outage capacity with \( N = 5 \) are computed based on [6] and [22], respectively. It can be observed that the derived theoretical results match perfectly with the simulation results, which validates the correctness of the theoretical results. In particular, the derived theoretical results are accurate enough with \( N = 5 \) no matter what the PS ratio is, which verifies our approximation again. Besides, the derived T2T outage capacity is more accurate than that in [10], which is one of the contributions of this paper. Another observation is that the achievable capacity increases first and then decreases with PS ratio, i.e., there only exists an optimal pair of symmetric PS ratio to maximize the achievable capacity.

Fig. 5 plots the optimal system outage capacity and PS ratio as a function of the distance of \( A - R \) link \( d_A \) with two PS schemes, respectively. The two PS schemes for comparison are the symmetric PS scheme and asymmetric PS scheme where both \( \lambda_A \) and \( \lambda_B \) are adjustable and the existing symmetric PS scheme with \( \lambda_A = \lambda_B \) as in [10]. The optimal system outage capacity and PS ratios for two PS schemes are obtained by maximizing the derived system outage capacity expression, as in (23). It is difficult to derive closed-form solutions for the problem (23) by using conventional solution methods. But we can use genetic algorithm-Based Algorithm to solve (23) (Please see Section 6 of [11] as an example). For convenience, we adopt the two-dimensional search algorithm to obtain the optimal PS ratios in Fig. 5 of the revised manuscript. It is assumed that the relay is located on the straight line between the both terminal nodes, i.e., \( d_A + d_B = 2 \). Note that here \( d_A = d_B \) results in the unequal statistic channel gains between the relay and the terminal nodes. One can see that the asymmetric PS scheme provides a higher capacity compared with the symmetric PS scheme when \( d_A \neq d_B \), while achieving the same capacity at \( d_A = d_B \). The reasons are as follows. The asymmetric PS scheme provides more flexibility than the symmetric PS scheme and the optimal PS ratios are determined by the statistic channel gains. If statistic channel gains are unequal, the asymmetric PS scheme makes two different optimal PS ratios and achieves a higher capacity; otherwise, the optimal PS ratios of the asymmetric PS scheme satisfy \( \lambda_A = \lambda_B \) and the PS ratio is equal to that of the symmetric PS scheme. One also can see that, in the asymmetric PS scheme, the relay harvests more (or less) energy with a larger (or smaller) PS ratio from the terminal node with a better (or worse) statistic channel to the relay, which achieves a trade-off between EH and information processing for the two terminal nodes.

Fig. 6 describes the relations of the optimal system outage
The optimal system outage capacity is to maximize
\[
\min \left\{ \gamma_{\text{th}} \right\} \text{subject to } \lambda_i, \gamma_{\text{th}} \geq \gamma_{\text{th}} \text{ for all } i.
\]
If \( \eta \) increases, \( \gamma_{\text{th}} \) will increase. In order to maximize the system outage capacity, \( \gamma_{\text{th}} \) should be improved, which can be realized by decreasing the optimal PS ratios. It can also be seen that the asymmetric PS scheme outperforms the symmetric one especially for a higher \( \eta \).

In Fig. 7, we plot system capacity versus static power allocation ratio \( \theta_A^2 \) with different relay locations for \( \lambda_A = \lambda_B = 0.5 \). It can be observed that the theoretical results match simulation results well, demonstrating the accuracy of our derived expression. Besides, there exists an optimal \( \theta_A^2 \) and the optimal \( \theta_A^2 \) is highly dependent on the relay location. To maximize the system capacity, the relay should assign less power to \( x_B \) when the relay is close to terminal A; otherwise, the relay should allocate more power to \( x_A \). Finally, Fig. 8 verifies the diversity analysis.

V. CONCLUSIONS

In this paper, we have studied the outage performance of three-step DF TWRNs with PS SWIPT. The expressions for T2T and system outage probabilities (and capacities) have been derived and verified by computer simulations. Some insights regarding the asymmetric PS scheme have been obtained. First, the optimal PS ratios are sensitive to the relay location and the energy conversion efficiency. In particular, the optimal \( \lambda_i \) decreases with the increase of the statistic channel gain of \( i \rightarrow R \) link. With the increase of the energy conversion efficiency, the optimal ratios, \( \lambda_A \) and \( \lambda_B \), show a downward trend. Second, the asymmetric PS scheme outperforms the symmetric one for the unequal statistic channel gains between the relay node and the terminal nodes. Third, the selection of power allocation ratio is highly related with the relay location.

This work can be extended by relaxing the assumption of perfect SIC, making our considered network close to practical scenario.

REFERENCES